

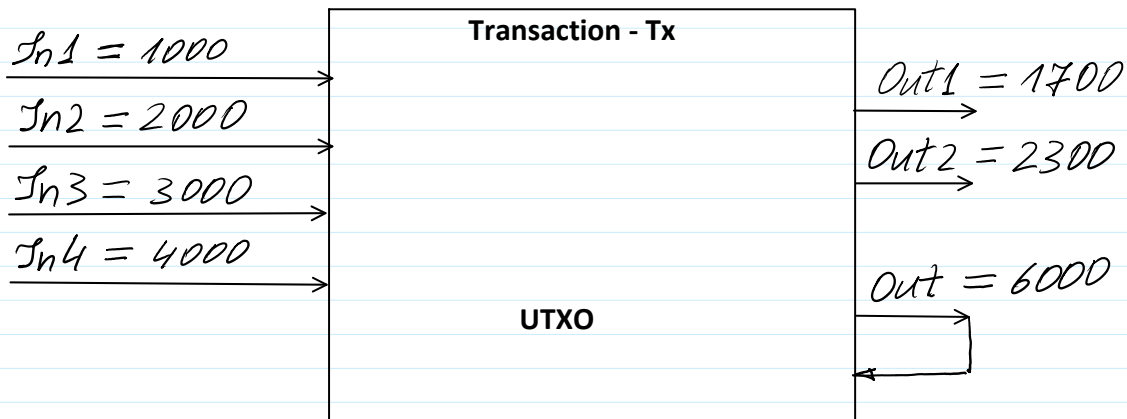
Interactive Exam Tasks:

1. Transaction creation according to template.
2. Public and private and public keys generation using $PP=(p, g)$; $p=264043379$, $g=2$.
3. Transaction signing using Schnorr signature.
4. Arbitrary 3 transactions including in the Merkle Tree.
5. To create a block and mine-validate created bloc.

Buhalterija - Book-Keeping, Accounting
Unspend Transaction Output - UTXO

No.	Pajamos-Incomes	Išlaidos-Expenses	Likutis-Balance
In1.	Client1: 1000 Sat		1000 Sat
In2.	Client2: 2000 Sat	Out1. Firm 5: 1700 Sat	1300 Sat
In3.	Client3: 3000 Sat	Out2. Firm 6: 2300 Sat	2000 Sat
In4.	Client4: 4000 Sat	Out3. Firm 7:	6000 Sat
Total	10 000 Sat	4000 Sat	6000 Sat

*Sum of Inputs =
 = Sum of Outputs
 Divisibility*



Public Parameters $PP = (p, g)$; $p=264043379$; $g=2$;

$A: PrK_A = x; PubK_A = a;$

$B: PrK_B = y; PubK_B = b;$

```
>> x=randi(2^27)
```

```
x = 100497451
```

```
>> a=mod_exp(g,x,p)
```

```
a = 91968695
```

```
>> y=randi(2^27)
```

```
y = 25824381
```

```
>> b=mod_exp(g,y,p)
```

```
b = 195335035
```

Schnorr Signature $\leftarrow M$ - message to be signed
 Signature is denoted by $\sigma = S(r, s)$

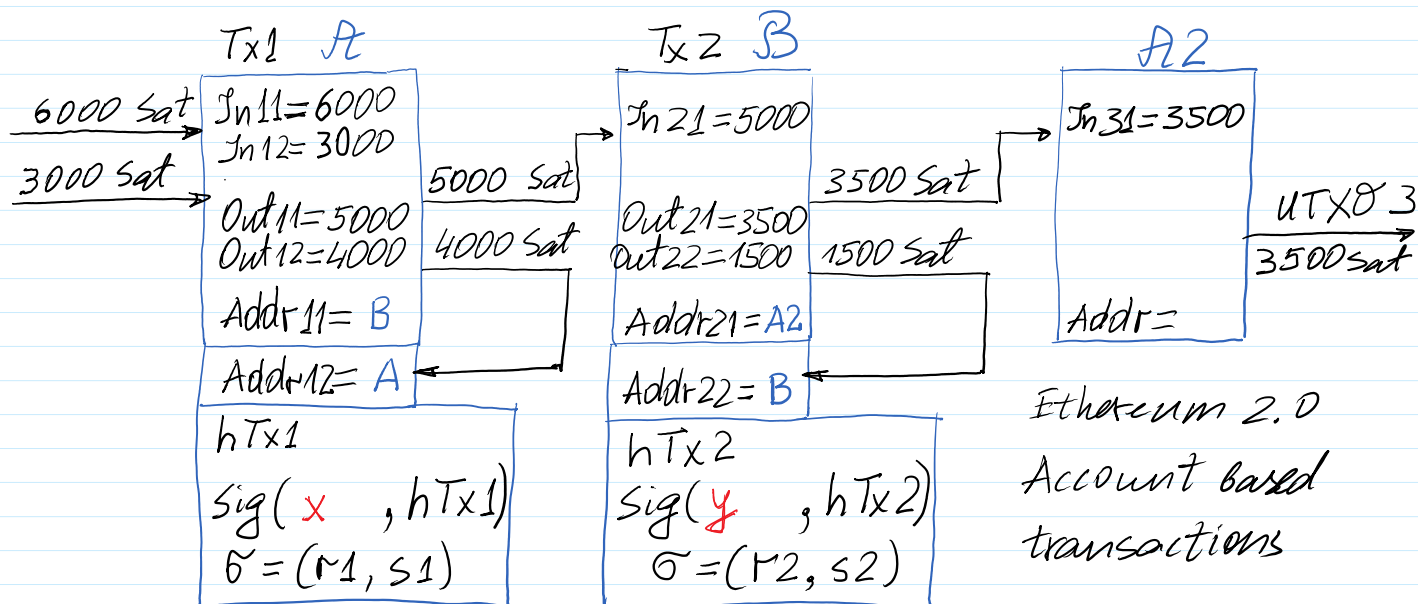
```
u=randi(p-1)
```

Verification of $\sigma = S(r, s)$ for $h = H(M || r)$

$u = \text{randi}(p-1)$
 $r = g^u \text{ mod } p$
 $h = H(M||r)$ % h is a decimal number
 $s = u + xh \text{ mod } (p-1)$

Verification of $\sigma = S = (r, s)$ for $h = H(M||r)$

$$g^s = r a^h \text{ mod } p$$



'Tx1: In11=6000 || In12=3000 || Out11=5000 || Out12=4000 || Rec1=A2 || Rec2=B'

$$h_{Tx1} = h_{28}(\downarrow)$$

$$\text{Sign}(PrK_A = x, h_{Tx1}) = s_1 = \tilde{\sigma}_1 = (r_1, s_1) \quad \% \text{ Schnorr Sig.}$$

'Tx2: In21=5000 || Out21=3500 || Out22=1500 || Rec1=A2 || Rec2=B'

$$h_{Tx2} = h_{28}(\downarrow)$$

$$\text{sign}(PrK_B = y, h_{Tx2}) = s_2 = \tilde{\sigma}_2 = (r_2, s_2) \quad \% \text{ Schnorr Sig.}$$

Transaction templates:

Tx1 = 'Tx1:In11=6000 || In12=3000 || Out11=5000 || Out12=4000 || Rec1=B || Rec2=A'

Tx2='Tx2:In21=5000 || Out21=3500 || Out22=1500 || Rec1=B2 || Rec2=B'

>> hTx1=h28('Tx1:In11=6000 || In12=3000 || Out11=5000 || Out12=4000 || Rec1=B || Rec2=A')

hTx1 = AFC73D8

>> hTx1=h28(Tx1)

hTx1 = AFC73D8

>> hTx1d=hd28(Tx1)

```
hTx1d = 184316888
```

```
>> hex2dec('AFC73D8')
```

```
ans = 184316888
```

```
>> hex2dec(hTx1)
```

```
ans = 184316888
```

```
>> hTx2=h28('Tx2:In21=5000||Out21=3500||Out22=1500||Rec1=B2||Rec2=B')
```

```
hTx2 = 43BC2D0
```

```
>> hTx2=h28(Tx2)
```

```
hTx2 = 43BC2D0
```

```
>> hTx2d=hd28(Tx2)
```

```
hTx2d = 71025360
```

```
>> hex2dec(hTx2)
```

```
ans = 71025360
```

Schnorr Signature $\leftarrow M$ - message to be signed in decimal format

Transaction Tx1 signing: $M=hTx1d = 184316888$

```
hTx1d=hd28(Tx1)
```

```
u1=randi(p-1)
```

```
r1=mod_exp(g,u1,p)
```

```
h1=hd28(hTx1d||r1) % h is a decimal number
```

```
s1=u1+x(h1) mod (p-1)
```

Signature for Tx1 computation $\sigma_1=S_1=(r_1, s_1)$

Signature for Tx2 computation $\sigma_2=S_2=(r_2, s_2)$

```
>> u1=randi(p-1)
```

```
u1 = 96927099
```

```
>> r1=mod_exp(g,u1,p)
```

```
r1 = 240239134
```

```
>> h1=hd28(hTx1d||r1)
```

```
h1 = 126174618
```

```
>> xh1=mod(x*h1,p-1)
```

```
xh1 = 319146
```

```
>> s1=mod(u1+xh1,p-1)
```

```
s1 = 97246245
```

$$g^s = r a^h \text{ mod } p$$

$$g^{s_1} = (r_1) a^{h_1} \text{ mod } p$$

```
>> g_s1=mod_exp(g,s1,p)
```

```
g_s1 = 168542564
```

```
>> a
```

```
a = 91968695
```

```
>> a_h1=mod_exp(a,h1,p)
a_h1 = 128485502
>> r1
r1 = 240239134
>> r1a_h1=mod(r1*a_h1,p)
r1a_h1 = 168542564
```

```
>> hTx1d=hex2dec('AFC73D8')
hTx1d = 184316888
```

%%% Signature **Sig1** Creation on **hTx1d**

Transaction Tx1 signing: **M=hTx1d**

```
u1=randi(p-1)
r1=mod_exp(g,u1,p)
h1=hd28('hTx1d||r') % h1 is a decimal number
xh1=mod(x*h1,p-1)
s1=mod(u1+xh1,p-1)
```

Sig1=(r1, s1)

Signature **Sig1** Verification

%%% Signature verification

```
>> hTx2d=hex2dec('43BC2D0')
hTx2d = 71025360
```

%%% Signature **Sig2** Creation on **hTx2d**

Transaction Tx2 signing: **M=hTx2d**

```
u2=randi(p-1)
r2=mod_exp(g,u1,p)
h2=hd28('hTx2d||r') % h1 is a decimal number
xh2=mod(y*h2,p-1)
s1=mod(u1+xh2,p-1)
```

Sig2=(r2, s2)

Signature **Sig2** Verification

%%% Signature verification

'Bl1: PrBl=... || Rh1=... || hTx1=... || hTx2=... || nonce=1000'

>> BI1M=h28('Bl1:PrBlh=0CAF06F || Rh=89B5B94 || hTx1=6707A9C || hTx2=C1169C4 || nonce=1000')
BI1M = D99C1E7

>> BI1M=h28('Bl1:PrBlh=0CAF06F || Rh=89B5B94 || hTx1=6707A9C || hTx2=C1169C4 || nonce=2018')
BI1M = 06B0772

h28: computes h-value of 28 bits

According to the difficulty target the adequate mined block must have leading hexadecimal 0 or 4 leading 0_b bits.

All available h-values with 28 bits length is equal to

>> 2²⁸ ans = 268435456

The number of adequate mined values is equal to any 28 - 4 = 24 bits values

>> 2²⁴ ans = 16777216

The probability to mine a block is the following:

$$Pr = \frac{\text{adequat number of values}}{\text{total number of values}} = \frac{2^{24}}{2^{28}} = 2^{-4} = \frac{1}{16}$$